

# eBay in the Sky: StrategyProof Wireless Spectrum Auctions

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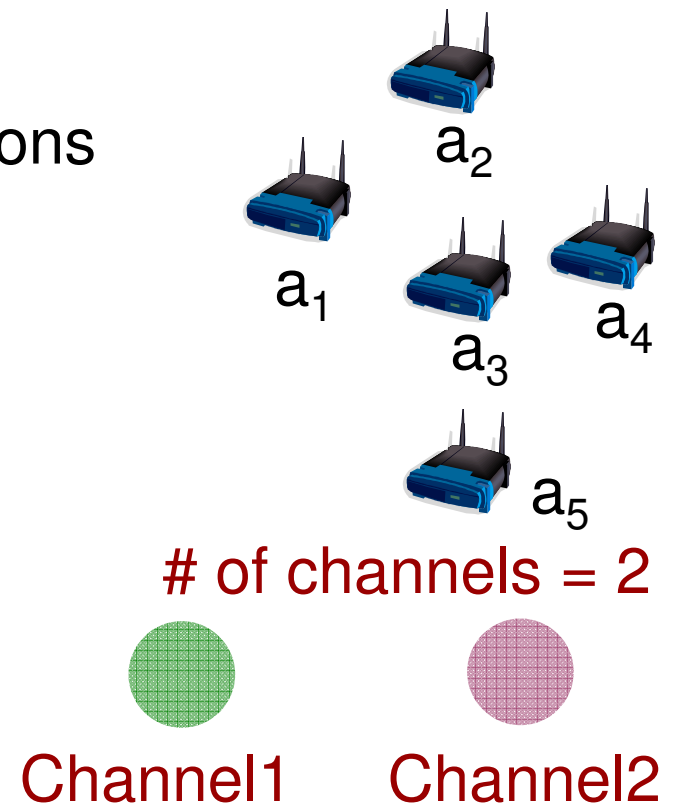
IUSTITIA  
(Goddess of Justice)



VERITAS  
(Goddess of Truth)

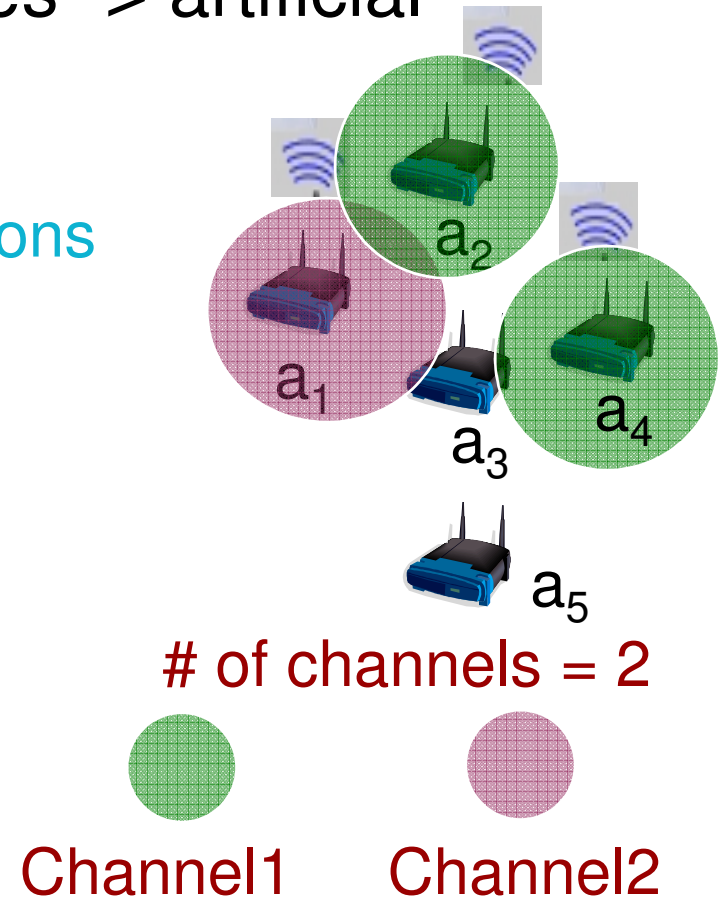
# Need for On-Demand Spectrum Auctions

- Explosion in the number of wireless devices
- FCC: static long term licenses -> artificial scarcity
- Solution:
  - Dynamic spectrum redistributions
  - Exploit spatial reusability



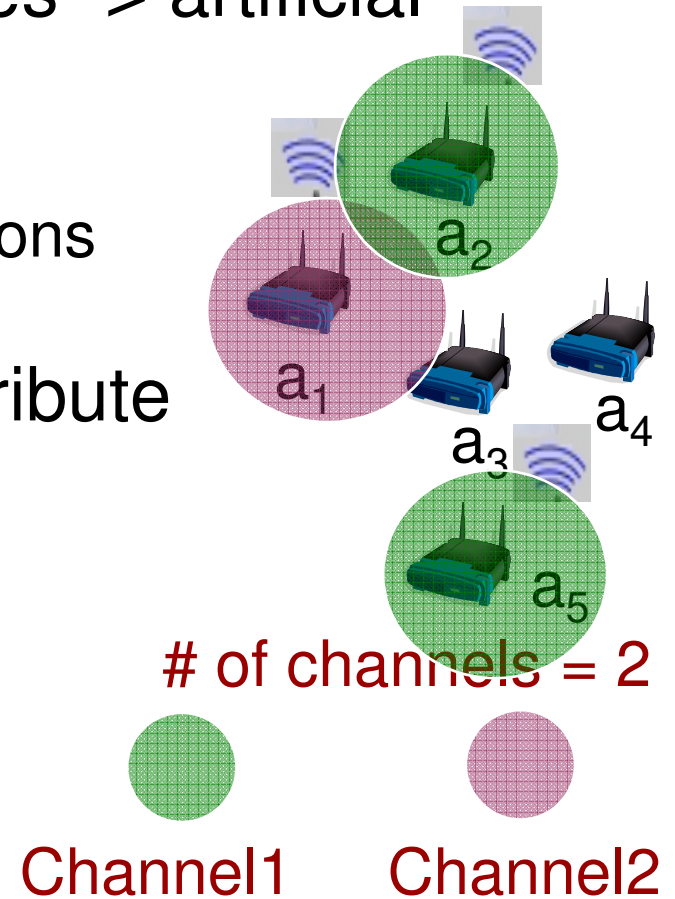
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- Explosion in the number of wireless devices
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- Solution:
  - Dynamic spectrum redistributions
  - Exploit spatial reusability
- Auctions widely used to distribute scarce resources
  - Fair and open
  - Economic Efficiency

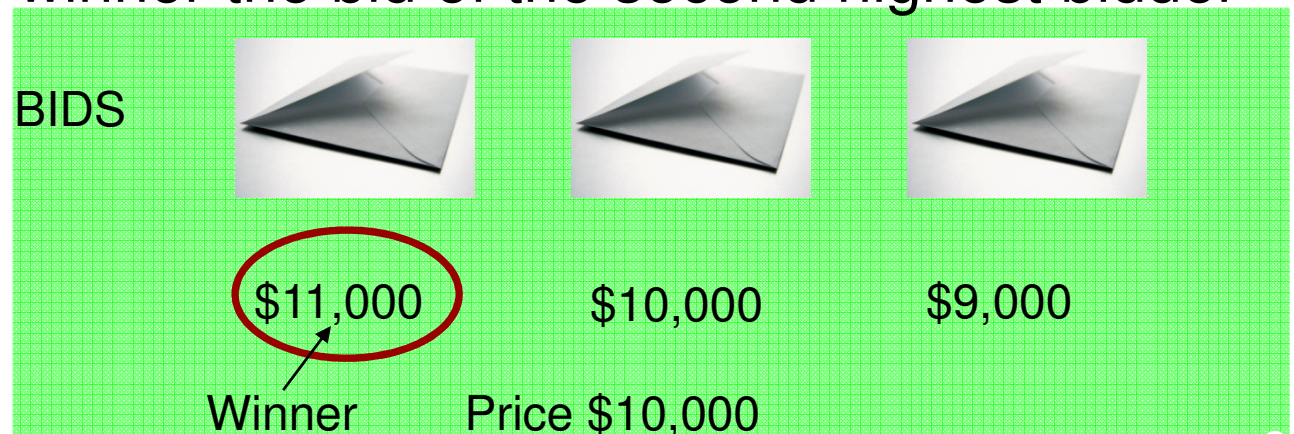


# Need for Truthful Mechanisms

- Selfish Bidders lead to challenges and overheads
  - Auction: Game among strategic players
  - Goal: Maximize individual utility
    - $Utility = (True\ Value - Price\ Paid)$
  - Players strategize over others
  - Counter productive for the auctioneer
- Truthful Auctions help overcome these challenges and overheads
  - Truthful auction: Every bidder maximizes its utility by bidding its true value

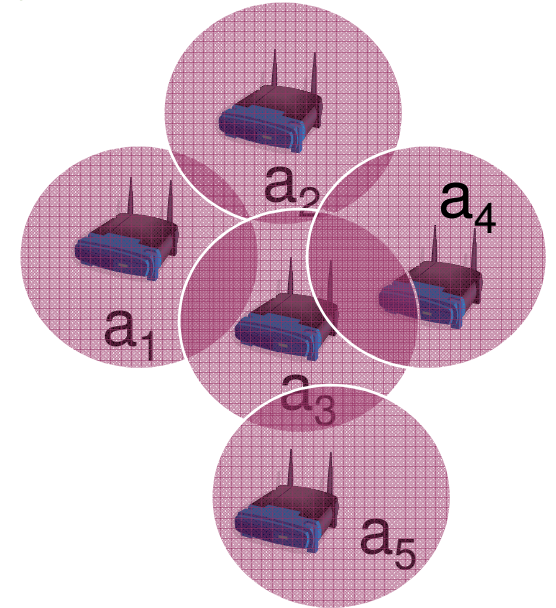
# Vickery: A Classical Truthful Auction

- Consider an auction for single item
- Vickery (Nobel Prize Winner)
  - Bidders submit bids in sealed envelopes
  - Auctioneer
    - Awards the item to the highest bidder
    - Charges winner the bid of the second highest bidder



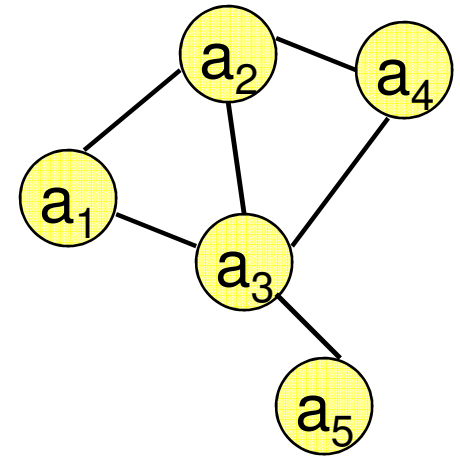
# Truthful Spectrum Auctions

- Items: Channels (k)
- Interference Graph
  - Nodes: Bidders
  - Edges: Interference Constraints



# Truthful Spectrum Auctions

- Items: Channels ( $k$ )
- Interference Graph
  - Nodes: Bidders
  - Edges: Interference Constraints
- Assumptions
  - Interference Graph is given
  - Static nodes
  - No collusion among bidders
  - Every bidder is bidding for one (any) channel

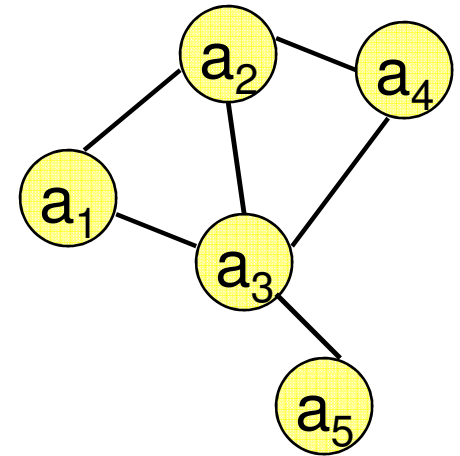


INTERFERENCE  
GRAPH



# Truthful Spectrum Auctions

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  - Node: Bidders
  - Edges: Interference Constraints
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INTERFERENCE  
GRAPH

# Possible Solution 1: Extending Vickery

- Algorithm (For allocating  $k$  channels)
  - Allocate channels to  $k$  highest bidders
  - Price: Bid of  $(k+1)$ th highest bidder

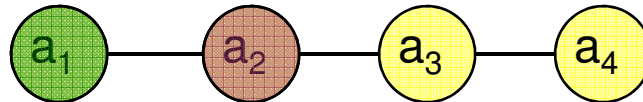
Bids

$b_1=5$

$b_2=4$

$b_3=1$

$b_4=2$



PRICE CHARGED : 2

# of channels = 2

- Inefficient spectrum utilization: spatial reuse not exploited

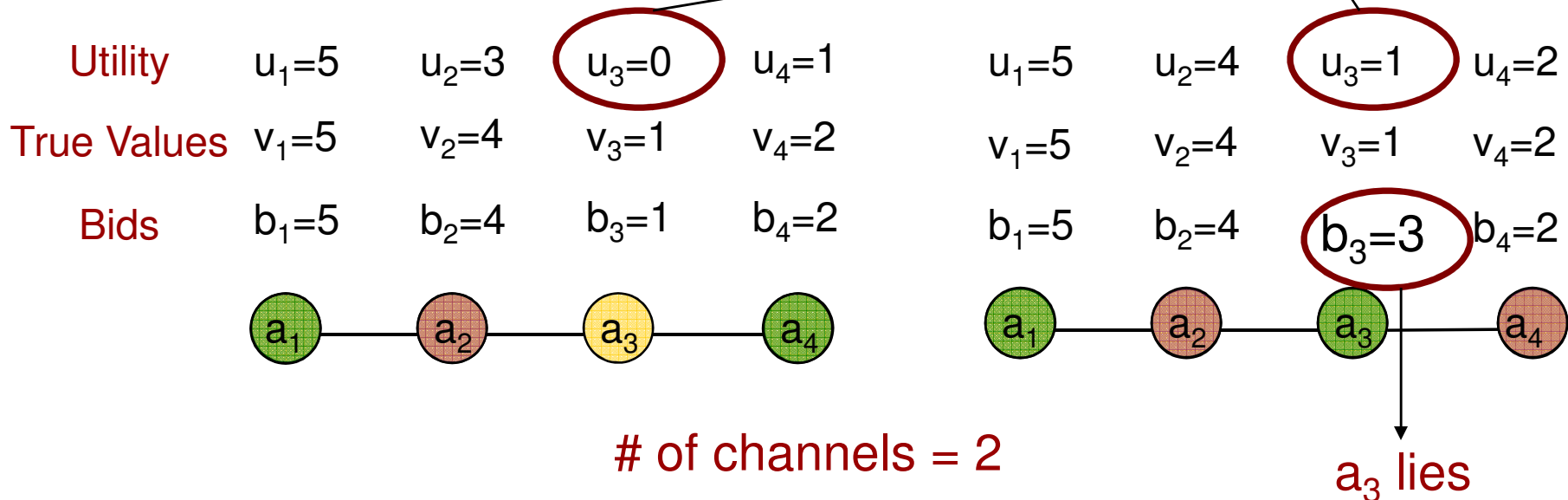
# Possible solution 2: VCG

- Optimal Spectrum utilization : NP-Hard
  - VCG is truthful but not polynomial time
- Relax Optimality constraint
  - **Pareto Optimal solution**: Cannot allocate any more channels without de-allocating at least one bidder
- Given an interference graph  $G$ , set of bids  $B$ , available number of channels  $k$ , design a truthful auction mechanism which run in **polynomial time**, results in **pareto optimal allocations** and has a **non-trivial pricing** scheme

# Possible Solution 3: Extending Secondary Price Auctions

- Sort and Greedily allocate channels
  - Allocate lowest available index
- Charge every winning bidder the bid of the highest unallocated neighbor

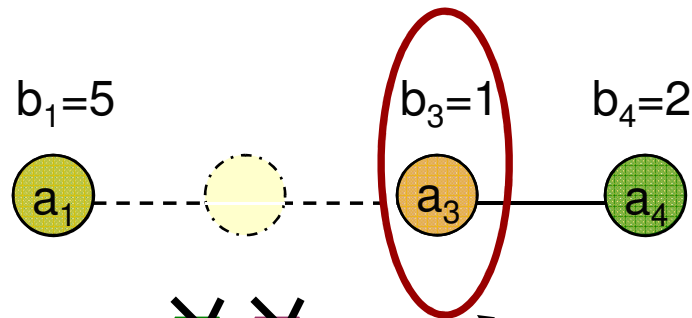
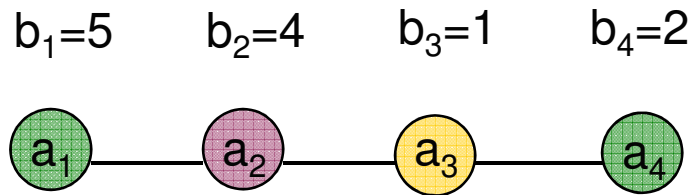
**VIOLATES TRUTHFULNESS !!!**



# Veritas: Truthful and Efficient Spectrum Auctions

- Greedy Allocation
  - Best known polynomial time channel allocation schemes are greedy
- Veritas-Pricing:
  - Charge every winner  $i$ , the bid of its **critical neighbor**  $C(i)$
  - Finding Critical Neighbor for  $i$ 
    - run allocations on  $\{B/b_i\}$  ( $B$ : set of bids)
    - *Critical Neighbor: The neighbor which makes the number of channels available for  $i$  go 0*

# Veritas Toy Example



Channels    
available for  $a_2$

Critical Neighbor for  $a_2$

# of channels = 2

# Proof of Veritas's truthfulness

- Theorem: Veritas spectrum auction is truthful, achieves pareto optimal allocations, and runs in polynomial time  $O(n^3k)$
- Proof sketch
  - **Critical Value**: Given a bid-set  $B$ , unique critical value exists for every allocated bidder.
  - **Monotonicity of allocations**: If a bidder bids greater its critical value, it is always allocated.
  - **Truthfulness**: If we charge every bidder by its critical value, no bidder has an incentive to lie.

# Simulations

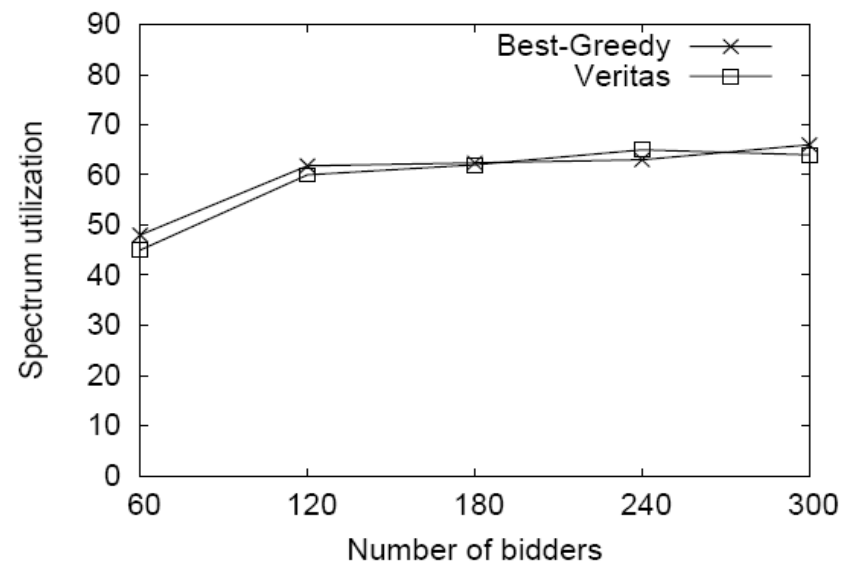
- Compare revenue and spectrum utilization of Veritas with other truthful and non-truthful designs.
- Synthetic Data
  - Nodes placed randomly in 1 X 1 square
  - Unit disk interference graph with radius = 0.1
  - Bids are randomly picked from the interval (0, 1]
- All results are averaged over multiple seeds



# Spectrum Utilization: Best-Greedy vs Veritas

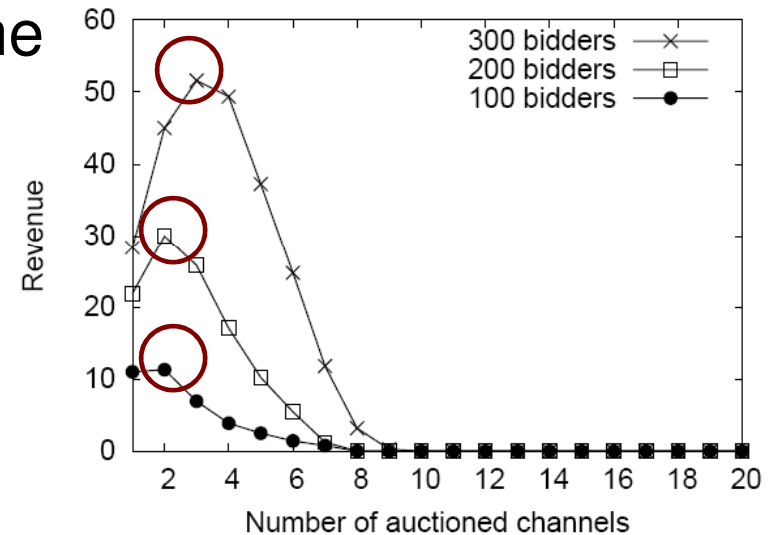
- Best Greedy: best known polynomial time spectrum allocation scheme (non-truthful)
- Veritas: Achieves truthfulness yet **comparable spectrum utilization**

# of channels = 8



# Veritas Revenue

- Revenue curve **not monotonically increasing** when # of channels is increased
  - Effect of truthful pricing scheme
  - Requires sufficient competition



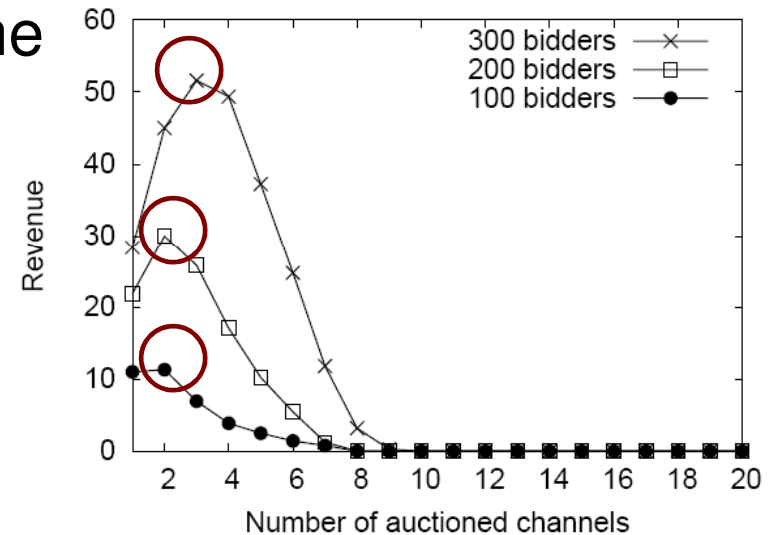
VERITAS ALGORITHM

# Veritas Revenue

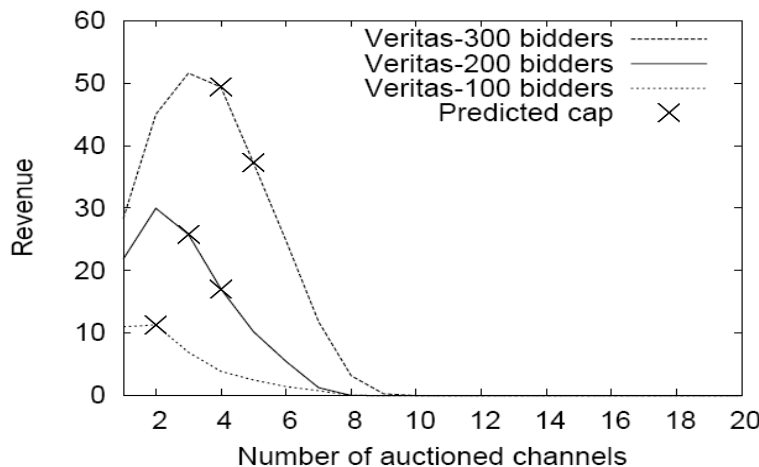
- Revenue curve **not monotonically increasing** when # of channels is increased

- Effect of truthful pricing scheme
- Requires sufficient competition

- Important to choose the right # of channels



VERITAS ALGORITHM



Heuristic for choosing number of channels

# Veritas Extensions

- **Objective functions:** Veritas allocation scheme can sort on broad class of functions of bids
  - Useful for choosing objective function
- **Bidding Formats:**
  - **Range Format:** Every bidder  $i$  specifies parameter  $d_i$ , and requests any number of channels in the range  $(0, d_i)$
  - **Contiguous Format:** Bidder requests the channels allocations to be contiguous

# Conclusion

- We propose Veritas: a **polynomial time** truthful mechanism for dynamic channel allocation (**pareto optimal allocations**)
- Related work: VERITAS differs from conventional spectrum allocation/auction designs by achieving both **truthfulness and spectrum efficiency (spatial reuse)**
- We show that the Veritas mechanism is **highly flexible** can be extended for many **objective functions** and **bidding formats**

# Thanks for listening

- Questions ???